# Addition Operation In Semigraphs<sup>\*</sup>

Jyoti Shetty<sup>†</sup>, Kadiyali Arathi Bhat<sup>‡§</sup>, Giliyaru Sudhakara<sup>¶</sup>

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#### Abstract

The addition of graphs is a well-studied operation on graphs which results in a new graph with more number of vertices and edges. The operation of adding vertices to graphs is different from the addition of graphs but also defines a new graph with more number of vertices and edges. In this article, the above two operations are extended to a generalization of graphs called semigraphs. The article also deals with the rank of a special incidence matrix of semigraphs resulting by the application of the above two operations.

### 1 Introduction

It is convenient to be able to express 'a given structure' in terms of smaller and simpler structures, which is possible only when we know the building blocks of the given structure. In graph theory many operations are defined which give raise to new graphs from the given graphs and the properties of new graphs in terms of their generators are studied. In this article, we extend some operations existing in the theory of graphs to semigraphs, a generalization of graphs.

The incidence matrix of a semigraph, as defined by E. Sampathkumar [9], does not represent the semigraph uniquely. Deshpande et al. [9], came up with a new definition of the incidence matrix of a semigraph, which represents the semigraph uniquely, but does not reveal the fact that  $(v_1, v_2, \ldots, v_n)$  and  $(v_n, v_{n-1}, \ldots, v_1)$  represent the same edge in a semigraph.

Authors in [10] make use of a property of binomial coefficient while defining the binomial incidence matrix, which not only represents a semigraph uniquely but also has the following property. The (i, j) entry of the matrix gives information about position of every vertex  $v_i$  incident on  $e_j$  from either end vertex on the edge  $e_j$  and also the size of the edge  $e_j$ .

### 2 Preliminaries

In this section, we give basics of semigraph [9] and readers are referred to [5] for all the elementary notations and definitions not described but used in this article.

**Definition 1** A semigraph G is a pair (V, E) where V is a nonempty set whose elements are called vertices of G, and E is a set of k-tuples of distinct vertices, called edges of G, for various  $k \ge 2$ , satisfying the following conditions.

1. Any two edges of G can have at most one vertex in common.

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<sup>&</sup>lt;sup>†</sup>Department of Mathematics, Manipal Institute of Technology Bengaluru, Manipal Academy of Higher Education, Manipal, Karnataka, India

<sup>&</sup>lt;sup>‡</sup>Department of Mathematics, Manipal Institute of Technology, Manipal Academy of Higher Education, Manipal, Karnataka 576104, India

<sup>&</sup>lt;sup>§</sup>Corresponding Author

 $<sup>\</sup>P$  Department of Mathematics, Manipal Institute of Technology, Manipal Academy of Higher Education, Manipal, Karnataka 576104, India

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- 2. Two edges  $(a_1, a_2, ..., a_p)$  and  $(b_1, b_2, ..., b_q)$  are said to be equal if and only if
  - p = q and
  - either  $a_i = b_i$  for  $1 \le i \le p$  or  $a_i = b_{p-i+1}$  for  $1 \le i \le p$ .

**Note 1** Let E be the set of vertices on the edge e. Then the size of E is called the size of the edge e and it is usually denoted by |E|.

Let G = (V, E) be a semigraph and let  $e = (u_1, u_2, \ldots, u_k)$  be an edge of G. Then  $u_1$  and  $u_k$  are called the end vertices and  $u_i, 2 \le i \le k-1$ , are called the mid vertices of e. Two vertices of G are adjacent if there is an edge containing both of them. An edge is said to be incident on every vertex on it. Two edges of G are adjacent if they have a vertex in common. Two vertices are consecutively adjacent if they are consecutive on the edge containing them. In a semigraph, an edge of size at least three is known as semiedge.

Like a graph, a semigraph G also has a geometric representation on plane. Vertices of G are represented either by dots or by small circles according to whether they are end vertices or mid vertices of the edge containing them and edges of G by curves passing through all the vertices on them. When a mid vertex v of an edge  $e_1$  is an end vertex of another edge, say  $e_2$ , then a small tangent is drawn to the circle representing vertex v where  $e_2$  meets v. A semigraph G and its representation are given in Example 1.

**Example 1** Let G = (V, E) be a semigraph with  $V(G) = \{u_1, u_2, \dots, u_8\}$  and

$$E(G) = \{e_1 = (u_1, u_2, u_3); e_2 = (u_3, u_4, u_5); e_3 = (u_2, u_8); e_4 = (u_5, u_6, u_7, u_8)\}$$

It can be represented as shown in Figure 1, and the edges  $e_1$ ,  $e_2$  and  $e_4$  are semiedges.



Figure 1: Semigraph G with 8 vertices and 4 edges

**Definition 2** A subedge of an edge  $e = (a_1, a_2, ..., a_m)$  is a p-tuple  $e' = (a_{i_1}, a_{i_2}, ..., a_{i_p})$ , induced by the vertices  $a_{i_1}, a_{i_2}, ..., a_{i_p}$ , where  $1 \le i_1 < i_2 < ... < i_p \le m$  or  $1 \le i_p < i_{p-1} < ... < i_1 \le m$ . A partial edge of e is a k-tuple  $e'(a_i, a_k) = (a_i, a_{i+1}, ..., a_{i+k-1})$ , where  $1 \le i \le m - k + 1$ . Here two consecutive vertices in e' are also consecutive vertices in e. Note that an edge is a subedge (partial edge) of itself, but a subedge (partial edge) is not an edge of G.

**Definition 3** A complete semigraph is a semigraph in which every two vertices are adjacent. In addition, if every vertex is an end vertex of some edge then it is called a strongly complete semigraph. For example, a semigraph which consists of a single edge of size  $n \ge 3$  is complete but not strongly complete and is denoted by  $E_n^c$ . The strongly complete semigraph on n vertices with one edge of size (n-1) and all other edges of size two is denoted by  $T_{n-1}^1$ .

**Definition 4** A semigraph G is said to be a zig-zag semigraph if  $V(G) = \{u_1, u_2, \ldots, u_k, u_{k+2}, \ldots, u_{2k-1}\}$ and

 $E(G) = \{(u_1, u_2, \dots, u_k), (u_1, u_{k+1}), (u_2, u_{k+1}), (u_2, u_{k+2}), \dots, (u_{k-1}, u_{2k-2}), (u_{k-1}, u_{2k-1}), (u_k, u_{2k-1})\}, denoted by Z_k^{k-1}.$ 

**Definition 5** A semigraph G' = (V', E') is a subsemigraph of a semigraph G = (V, E) if  $V' \subseteq V$  and  $E' \subseteq E$ .

**Definition 6** Let  $G_1$  and  $G_2$  be two graphs with disjoint vertex sets  $V_1$  and  $V_2$  and edge sets  $E_1$  and  $E_2$ , respectively. Their union denoted by  $G_1 \cup G_2$  has vertex set  $V = V_1 \cup V_2$  and edge set  $E = E_1 \cup E_2$ . Their join (addition) denoted by  $G_1 + G_2$  consists of  $G = G_1 \cup G_2$  and all edges joining  $V_1$  with  $V_2$ .

Now we define the binomial incidence matrix of a semigraph.

**Definition 7** ([10]) Let G = (V, E) be a semigraph with  $V(G) = \{u_1, u_2, \ldots, u_n\}$  and  $E(G) = \{e_1, e_2, \ldots, e_m\}$ . Let size of the edge  $e_j$  be  $n_j + 1$ ,  $1 \leq j \leq m$ . The binomial incidence matrix of G, denoted by  $\mathscr{B}(G)$ , is a  $n \times m$  matrix, whose rows are indexed by the vertex set and columns are indexed by the edge set of G. The column corresponding to  $e_j$  in the binomial incidence matrix consists of entries  $0, {}^{n_j}C_0, \ldots, {}^{n_j}C_{n_j}$ , where nonzero entries correspond to the vertices on the edge. The entries  ${}^{n_j}C_0$  and  ${}^{n_j}C_{n_j}$  correspond to the end vertices of the edge  $e_j$ . The (i, j) entry of  $\mathscr{B}(G)$  is given by

$$b_{ij} = \begin{cases} \begin{array}{cc} {}^{n_j}C_r & \text{if vertex } u_i \text{ and edge } e_j \text{ are incident and } u_i \text{ is the } r^{th} \text{ vertex} \\ & \text{from the end vertex of } e_j \text{ with entry } {}^{n_j}C_0, \ 0 \le r \le n_j \\ 0 & \text{if vertex } u_i \text{ and edge } e_j \text{ are not incident on each other.} \end{cases} \end{cases}$$

**Example 2** The binomial incidence matrix  $\mathscr{B}(G)$  of semigraph G as shown in Figure 1 is given by

$$\mathscr{B}(\mathbf{G}) = \begin{array}{ccccccc} & e_1 & e_2 & e_3 & e_4 \\ & u_1 \\ & u_2 \\ & u_3 \\ & u_4 \\ & u_5 \\ & u_6 \\ & u_7 \\ & u_8 \end{array} \begin{pmatrix} {}^2C_0 & 0 & 0 & 0 \\ {}^2C_1 & 0 & {}^1C_0 & 0 \\ {}^2C_2 & {}^2C_0 & 0 & 0 \\ 0 & {}^2C_2 & 0 & {}^3C_0 \\ 0 & 0 & 0 & {}^3C_1 \\ 0 & 0 & 0 & {}^3C_2 \\ 0 & 0 & {}^1C_1 & {}^3C_3 \end{pmatrix}.$$

We refer the interested readers to the articles [2, 3, 4, 6, 7, 8] for the other important studies in semigraphs.

## 3 Number of Subedges and Number of Partial Edges in a Semigraph

The semigraph  $E_n^c$ ,  $n \geq 3$  seems to be the simplest semigraph which is not a graph.

**Example 3** The semigraph  $E_3^c$  is as shown in Figure 2. Subedges of the semigraph G are  $(v_1, v_2, v_3)$ ,  $(v_1, v_2)$ ,

Figure 2: Semigraph  $G = E_3^c$ 

 $(v_1, v_3)$  and  $(v_2, v_3)$  out of which  $(v_1, v_2, v_3)$ ,  $(v_1, v_2)$ , and  $(v_2, v_3)$ , are partial edges.

It is quite interesting to count the number of subedges and the number of partial edges in a given semigraph and the task can be simplified by knowing the number of subedges and partial edges in an edge of size n. Shetty et al.

**Theorem 1** The number of subedges in an edge of size n is given by  $2^n - n - 1$  and that of partial edges is  ${}^{n}C_2$ .

**Proof.** Let  $e = (v_1, v_2, \ldots, v_n)$  be an edge of size n. With  $v_1$  as starting vertex, we have (n-1) partial edges namely  $[(v_1, v_2), (v_1, v_2, v_3) \ldots (v_1, v_2, \ldots, v_n)]$ . Similarly, with  $v_2$  as starting vertex there are (n-2) partial edges.

Finally, with  $v_{n-1}$  as starting vertex, we have one partial edge. Hence, number of partial edges in e is,

$$1 + 2 + 3 + \ldots + n - 1 = \frac{n(n-1)}{2}$$
.

With  $v_1$  as starting vertex, there are  ${}^{n-1}C_1$  subedges of size 2 namely  $[(v_1, v_2), (v_1, v_3) \dots (v_1, v_n)]$ . From  $v_1$  there are  ${}^{n-1}C_2$  subedges of size 3 and so on. With  $v_1$  as starting vertex there is one subedge of size n. Hence the total number of subedges in e with  $v_1$  as starting vertex is,

$$^{n-1}C_1 + {}^{n-1}C_2 + \ldots + {}^{n-1}C_{n-1} = 2^{n-1} - 1.$$

Out of  $2^{n-1} - 1$  subedges with  $v_1$  as starting vertex, there are (n-1) partial edges.

Similarly, with  $v_2$  as starting vertex, there are  $2^{n-2} - 1$  subedges in e, out of which (n-2) are partial edges and so on. With  $v_{n-1}$  as starting vertex, there is only one subedge which is also a partial edge. Hence, total number of subedges in the edge e which is of size n is,

$$(2^{n-1}-1) + (2^{n-2}-1) + \ldots + (2^1-1) = 2^n - n - 1.$$

**Corollary 1** The number of subedges of size 2, in an edge of size n is given by  ${}^{n}C_{2}$  and that of partial edges of size 2, is n - 1.

**Proof.** Let  $e = (v_1, v_2, \ldots, v_n)$  be an edge of size n. With  $v_1$  as starting vertex, there are  ${}^{n-1}C_1$  subedges of size 2, out of which  $(v_1, v_2)$  is a partial edge.

From  $v_2$ , there are  $^{n-2}C_1$  subedges of size 2, out of which  $(v_2, v_3)$  is a partial edge.

:

With  $v_{n-1}$  as starting vertex, there is only one subedge of size 2, which is also a partial edge. Hence, total subedges of size 2 in the edge e of size n is,

$$1 + 2 + \ldots + {^{n-2}C_1} + {^{n-1}C_1} = {^nC_2}$$

Each vertex with its immediate next vertex forms a partial edge each of size 2. Those are  $(v_1, v_2)$ ,  $(v_2, v_3)$ , ...,  $(v_{n-1}, v_n)$ . Hence total partial edges of size 2 are, n-1.

### 4 Addition of a Complete Graph and Path Graph to Semigraph

In this section, we discuss about adding a complete graph and path graph to semigraphs. We have also explored the rank of the binomial incidence matrix of the resulting semigraphs.

**Definition 8**  $T_r^{K_{n-r}}$  is a complete semigraph on n vertices which contains a semiedge of size  $r \ge 3$  and a complete graph  $K_{n-r}$  as subsemigraphs, all the vertices of  $K_{n-r}$  are made adjacent to all the vertices on the semiedge by edges of size two. Clearly, number of edges in  $T_r^{K_{n-r}}$  is given by  $1 + {}^{n-r}C_2 + (n-r)r$ .

**Example 4** The semigraph  $T_3^{K_4}$  as shown in Figure 3 is a complete semigraph on 7 vertices which contains a semiedge of size 3 and a complete graph  $K_4$  as subsemigraphs.



Figure 3: Semigraph  $T_3^{K_4}$ 



Figure 4: Semigraph  $T_3^{P_4}$ 

**Definition 9**  $T_r^{P_{n-r}}$  is a semigraph on n vertices which contains a semiedge of size  $r \ge 3$  and a path graph  $P_{n-r}$  as subsemigraphs, all the vertices of  $P_{n-r}$  are made adjacent to all the vertices on the semiedge by edges of size two. Clearly, number of edges in  $T_r^{P_{n-r}}$  is given by (n-r)r+1.

**Example 5** The semigraph  $T_3^{P_4}$  as shown in Figure 4 is a semigraph on 7 vertices which contains a semiedge of size 3 and a path graph  $P_4$  as subsemigraphs.

Note 2 Clearly,  $T_r^{K_{n-r}} = E_r^c + K_{n-r}$  and  $T_r^{P_{n-r}} = E_r^c + P_{n-r}$ , when '+' represents addition of graphs.

We have the following theorem.

**Theorem 2** Let H be any connected graph on (n-r) vertices and G be the semigraph given by  $G = E_r^c + H$ . Then, rank of  $\mathscr{B}(G)$  is n.

**Proof.** Let  $v_1, v_2, \ldots, v_r$  be vertices of  $E_r^c$  and  $v_{r+1}, v_{r+2}, \ldots, v_n$  be the vertices of H. The first column of  $\mathscr{B}(G)$  has non zero entries  $r^{-1}C_0, r^{-1}C_1, \ldots, r^{-1}C_{r-1}$ . Assume that they appear in that order in the first r rows.

Suppose,  $X \in \mathbb{R}^n$  be such that  $X'\mathscr{B}(G) = 0$ . Since, for every j with  $1 \leq j \leq r$ , the vertex  $v_j$  is adjacent to  $v_{r+i}$ , for all  $i, 1 \leq i \leq n-r, x_j + x_{r+i} = 0$ , i.e  $x_j = -x_{r+i}, 1 \leq j \leq r$ . We also have,

Thus X is the zero vector. The dimension of the left null space of  $\mathscr{B}(G)$  is zero. Hence the result.

**Corollary 2** Let  $H = E_{n-r}^c$  and  $G = E_r^c + H = E_r^c + E_{n-r}^c = T_r^H$  is a complete semigraph on n vertices with r(n-r) + 2 edges. Then rank  $\mathscr{B}(T_r^H) = n$ .

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**Proof.** The proof follows from Theorem 2.  $\blacksquare$ 

- Note 3 1. If  $H = K_{n-r}$  then  $T_r^H = T_r^{K_{n-r}}$  which is a complete semigraph on n vertices having  $1 + n^{-r}C_2 + (n-r)r$  edges.
  - 2. If H is the complete semigraph with k edges, then  $T_r^H = E_r^c + H$  has n vertices and 1 + r(n-r) + k edges with the rank of  $\mathscr{B}(T_r^H)$  equal to n.
  - 3. Both the semigraphs  $T_r^{K_{n-r}}$  and  $T_r^{P_{n-r}}$  have the same rank n.
  - 4. Generalizing Corollary 2, if  $r_1, r_2, \ldots, r_k$  are positive integers greater than or equal to 2 with  $r_1 + r_2 + \ldots + r_k = n$ , then  $G = E_{r_1}^c + E_{r_2}^c + \ldots + E_{r_k}^c$  is a semigraph obtained by joining every vertex of  $E_{r_j}^c$  to every vertex which is not on  $E_{r_j}^c$ ,  $1 \le j \le k$ . This results in a complete semigraph on n vertices and  $(r_1 \times r_2 \times \ldots \times r_k) + k$  edges with the rank of  $\mathscr{B}(G)$  equal to n.

### 5 Rank of Binomial Incidence Matrix of a Semigraph after Addition of Vertices

In [1], Jean H. Bevis et al. have defined the addition of vertices to a given graph and discussed its effect on the rank of the resulting graph. They have defined the following operations on an undirected graph H.

- 1. Let *H* be an undirected graph, *u* be a vertex of *H* and *v* be a vertex not in V(H). Then  $H_1 = H \oplus_u v$  is the graph obtained by adding the vertex *v* and the edge between *u* and *v*, to H.
- 2. Let *H* be an undirected graph with  $V(H) = \{u_1, u_2, \ldots, u_n\}$ , *v* be a vertex not in V(H) and *X* be a 0-1 vector with *n* components. Then  $H_2 = H \oplus_X v$  is the graph obtained by adding the vertex *v* and making it adjacent with  $u_i$  only if  $i^{th}$  component of *X* is  $1, 1 \leq i \leq n$ .

Motivated by the above, we have defined the following operations on semigraphs. Let G be a semigraph with  $V(G) = \{u_1, u_2, \ldots, u_n\}.$ 

1. Let  $v \notin V(G)$  and let X be a 0-1 vector with n components. Then  $G \oplus_X v$  is the semigraph obtained by adding vertex v to G and making it adjacent to  $u_i$  by an edge of size 2, only if  $i^{th}$  component of X is 1,  $1 \leq i \leq n$ .

If  $X = \mathbf{1}$ , a vector in which all components are equal to 1, then  $G \oplus_X v = G \oplus_\mathbf{1} v$  is the semigraph obtained by making v adjacent to all vertices of G by the edges  $(u_i, v)$  of size 2,  $\forall i, 1 \leq i \leq n$ .

**Example 6** The complete semigraph  $T_5^{K_1} = [E_5^c \oplus_1 u_6]$  is as shown in Figure 5.



Figure 5: Complete semigraph  $T_5^{K_1}$ 

**Note 4** Let  $\{u_1, u_2, \ldots, u_n\}$  be the set of vertices of the complete semigraph  $E_n^c$  and  $G = E_n^c \oplus_1 u_{n+1} = E_n^c + K_1 = T_n^{K_1}$ . Then from Theorem 2,  $\mathscr{B}(G)$  is a full rank matrix.

2. Let  $u_{n+1}, u_{n+2}, \ldots, u_{2n-1}$  be vertices such that  $u_{n+i} \notin V(G), \forall i, 1 \le i \le n-1$ . Then

$$\left[G \oplus_{\{u_{i}, u_{i+1}\}}' u_{n+i}\right]_{i=1}^{n-1}$$

is a semigraph obtained by adding the edges of size 2 joining  $u_{n+i}$  and  $u_i$ ,  $i, 1 \le i \le n-1$ .

**Example 7** The Zig-zag semigraph,  $Z_6^5 = \left[E_6^c \oplus_{\{u_i, u_{i+1}\}}^{\prime} u_{n+i}\right]_{i=1}^5$  is as shown in Figure 6.



Figure 6: Zig-zag semigraph  $Z_6^5$ 

**Note 5** Let  $G = E_n^c$  be a semigraph with vertex set  $\{u_1, u_2, \ldots, u_n\}$ , and

$$G' = \left[ G \oplus_{\{u_i, u_{i+1}\}}' u_{n+i} \right]_{i=1}^{n-1} = Z_n^{n-1}.$$

Then  $\mathscr{B}(G')$  is a full rank matrix.

3. Let  $u_{n+1}, u_{n+2}, \ldots, u_{n+N}$  be vertices such that  $u_{n+i} \notin V(G)$ ,  $\forall i, 1 \le i \le N$ . Then  $\left[G \oplus_{\{u_1, u_n\}}^{''} u_{n+i}\right]_{i=1}^{N}$  is a semigraph obtained by making the vertices  $u_{n+1}, u_{n+2}, \ldots, u_{n+N}$  adjacent to both  $u_1$  and  $u_n$  by edges of size 2.

**Example 8** The semigraph,  $G' = \left[E_5^c \oplus_{\{u_1,u_5\}}^{''} u_{5+i}\right]_{i=1}^6$  is as shown in Figure 7.



Figure 7: Semigraph  $G' = \left[ E_5^c \oplus_{\{u_1, u_5\}}'' u_{5+i} \right]_{i=1}^6$ 

4. Let  $u_{n+1}, u_{n+2}, \ldots, u_{n+n}$  be vertices of  $H = E_n^c$  such that  $u_{n+i} \notin V(G)$ ,  $\forall i, 1 \le i \le n$ . Then  $\left[G \oplus_{u_i}^{''} u_{n+i}\right]_{i=1}^n$  is a semigraph obtained by making the vertices  $u_{n+1}, u_{n+2}, \ldots, u_{n+n}$  adjacent to the corresponding vertex of H by edges of size 2.



Figure 8: Semigraph  $G' = \left[ E_4^c \oplus_{u_i}^{'''} u_{n+i} \right]_{i=1}^4$ 

**Example 9** The semigraph,  $G' = \left[E_4^c \oplus_{u_i}^{\prime\prime\prime} u_{n+i}\right]_{i=1}^4$  is as shown in Figure 8.

**Theorem 3** Let  $G = E_n^c$  with  $V(G) = \{u_1, u_2, \ldots, u_n\}$  and let  $G' = \left[G \oplus_{\{u_1, u_n\}}^{''} u_{n+i}\right]_{i=1}^N$  where N is any positive integer. Then rank of  $\mathscr{B}(G')$  is N+2.

**Proof.** G' has n + N vertices  $u_1, u_2, \ldots, u_n, u_{n+1}, u_{n+2}, \ldots, u_{n+N}$  and 2N + 1 edges given by,  $e_1 = (u_1, u_2, \ldots, u_n), e_{i+1} = (u_1, u_{n+i})$  and  $e'_{i+1} = (u_n, u_{n+i}), 1 \le i \le N$ .

Then,  $\mathscr{B}(G')$  is of the form

	$e_1$	$e_2$	$e_3$	 $e_{N+1}$	e'	$'_{2}$	$e'_3$		$e'_{N+1}$
$u_1$	$\binom{n-1}{C_0}$	1	1	 1					
$u_2$	$^{n-1}C_1$		$O_{n-1 \times n-1}$					$O_{n-1 \times n-1}$	
$u_3$	${}^{n-1}C_2$		$O_{n-1 \times n-1}$						
•									
$u_n$	$^{n-1}C_{n-1}$				1	1	1		1
$u_{n+1}$	0				1				
$u_{n+2}$	0								
•			$I_{N \times N}$		1		$I_{N \times N}$		
$u_{n+N}$	0				1				)

where  $I_{N \times N}$  is the identity matrix of order N and  $O_{n-1 \times n-1}$  is the zero matrices of order n-1. One can easily observe that rank of  $\mathscr{B}(G') = N + 2$ .

**Corollary 3** Let  $G = E_n^c$  be a semigraph with vertex set  $\{u_1, u_2, \ldots, u_n\}$ , and let  $G' = \left[G \oplus_{\{u_1, u_n\}}^{''} u_{n+i}\right]_{i=1}^{N}$ , where N = n - 1 and det  $\mathscr{B}(G')$  is equal to zero. **Proof.** When N = n - 1, the matrix  $\mathscr{B}(G')$  becomes a square matrix of order (2n - 1). From Theorem 3, rank of  $\mathscr{B}(G') = (n - 1) + 2 = n + 1$ . Hence G' is not a full rank matrix and det  $\mathscr{B}(G') = 0$ .

**Theorem 4** Let  $G = E_n^c$  be a semigraph with vertex set  $\{u_1, u_2, \ldots, u_n\}$  and let  $u_{n+1}, u_{n+2}, \ldots, u_{n+n}$  be vertices of  $H = E_n^c$ . Let  $G' = \left[G \oplus_{u_i}^{'''} u_{n+i}\right]_{i=1}^n$ . Then the rank of  $\mathscr{B}(G')$  is m-1 where m is the number of edges in G'.

**Proof.** The number of vertices and edges in G' is 2n and n+2, respectively. Clearly,  $2n > n+2 \forall$ ,  $n \ge 3$ . The binomial incidence matrix of G' is of the form

If  $X = (x_1, x_2, \dots, x_{n+2})^T$  is a vector in the right null space  $\mathscr{B}(G')$  then  $\mathscr{B}X = 0$ . This implies that,

Hence, there exits only one vector in the right null space of  $\mathscr{B}X = 0$  i.e  $(1, 1, -^{n-1}C_0, \ldots, -^{n-1}C_{n-1})$ . Therefore, rank of  $\mathscr{B}(G') = (n+2) - 1 = m-1$ .

**Theorem 5** Let G = (V, E) a connected semigraph with exactly one semiedge. If G contains  $T_{n-1}^1$  or  $Z_n^{n-1}$  as its subsemigraph then rank of  $\mathscr{B}(G)$  is equal to |V(G)|.

**Proof.** Let G be a semigraph with p vertices and m edges among which only one edge is a semiedge and let  $T_{n-1}^1$  be a subsemigraph of G. We note that,  $n \le p \le m$ .

Let  $\mathscr{B}(G)$  be the binomial incidence matrix of G Then it is of the form,

$$\mathscr{B}(\mathbf{G}) = \begin{bmatrix} \mathscr{B}(T_{n-1}^1)_{n \times n} \\ O_{p-n \times n} \end{bmatrix} F ],$$

where  $\mathscr{B}(T_{n-1}^1)$  is a square matrix of order n with rows as parts of first n rows of  $\mathscr{B}(G)$  and column corresponding to the edges  $e_1, e_2, \ldots, e_n$  which are first n edges of G. And,  $O_{p-n\times n}$  is a zero matrix of order  $(p-n)\times n$  and F is  $p\times (m-n)$  matrix with columns corresponding to edges  $e_{n+1}, e_{n+2}, \ldots, e_m$  which are of size two.

Suppose,  $X \in \mathbb{R}^p$  be such that  $X'\mathscr{B}(G) = 0$ . Then,  $x_i + x_j = 0$  whenever the vertices  $u_i$  and  $u_j$  are adjacent by an edge of size two. Therefore,  $x_i + x_n = 0$  that is  $x_i = -x_n$ ,  $\forall i, 1 \le i \le n-1$ . We have from first column of  $\mathscr{B}(G)$ ,

$$x_1^{n-2}C_0 + x_2^{n-2}C_1 + \dots + x_{n-1}^{n-2}C_{n-2} = 0,$$
  
i.e  $-x_n \{ {}^{n-2}C_0 + {}^{n-2}C_1 + \dots + {}^{n-2}C_{n-2} \} = 0$   
 $\implies x_n = 0 \implies x_i = 0, \ \forall i, \ 1 \le i \le n.$ 

Now, consider the partition  $\{V_1, V_2\}$  of V(G) where  $V_1$  is the vertex set of  $T'_{n-1}$  and  $V_2 = V(G) \setminus V_1$ . For an edge  $e_i$ , i > n with end vertices  $u_s$  and  $u_t$ , since  $e_i$  is of size two, we have  $x_s + x_t = 0$ . If  $u_s \in V_1$  and  $u_t \in V_2$ , then  $s \leq n$  and t > n and  $x_s = 0$ , which implies  $x_t = 0$ . Suppose that  $u_s$  and  $u_t$  both are in  $V_2$ . Then, since G is connected, there exists a path between  $u_s$  and  $u_t$  which consists only of edges of size two and which passes through at least one vertex in  $V_1$ . This implies that  $x_s = 0$  and  $x_t = 0$ . Hence,  $x_i = 0 \quad \forall i, 1 \leq i \leq p$ . Thus, null space of  $\mathscr{B}(G)$  is of dimension zero and rank of  $\mathscr{B}(G)$  is equal to |V(G)|.

Continuing in similar lines, one can prove that rank of  $\mathscr{B}(G) = |V(G)|$  when G is a connected semigraph which has only one semiedge and has  $Z_n^{n-1}$  as a sub semigraph.

**Applications** In communication networks, not all nodes are equally important. Some of them may just receive and transmit the data whereas, some others may be used to process the data. So, in the graph representing the communication networks, all the vertices need not have the same importance, such situations could be better represented by a graph model called semigraph. In particular,  $E_n^c$  is an edge with *n* vertices on the same edge with exactly two end vertices and the remaining are mid vertices, where any two vertices are mutually adjacent to each other. Hence, in communication networks,  $E_n^c$  represents the situation where the communication takes place between every pair of vertices and only two of them which are placed at the end are of more importance.

In theoretical computer science, communication complexity studies the amount of communication required to solve a problem when the input to the problem is distributed among two or more parties. In the field of communication complexity, the rank of the communication matrix of a function gives bounds on the amount of communication needed for two or more parties to compute the function.

In this article, we have studied the rank of the binomial incidence matrix of the semigraph which represents the semigraph uniquely, which inturn gives a better picture of communication complexity when  $E_n^c$  is part of the communication network.

**Conclusions** We have studied the structure of a semigraph in terms of its substructures and also change in the structure of semigraph by adding graphs and vertices to a specific semigraph. Also the effect on the rank of the binomial incidence matrix of a semigraph have been discussed. Majority of the semigraphs resulted from addition operation in this article are of full rank.

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