

# Numerical Analysis I

## Solutions of Equations in One Variable

Instructor: Wei-Cheng Wang<sup>1</sup>

Department of Mathematics  
National TsingHua University

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<sup>1</sup>These slides are based on Prof. Tsung-Ming Huang(NTNU)'s original slides

# Outline

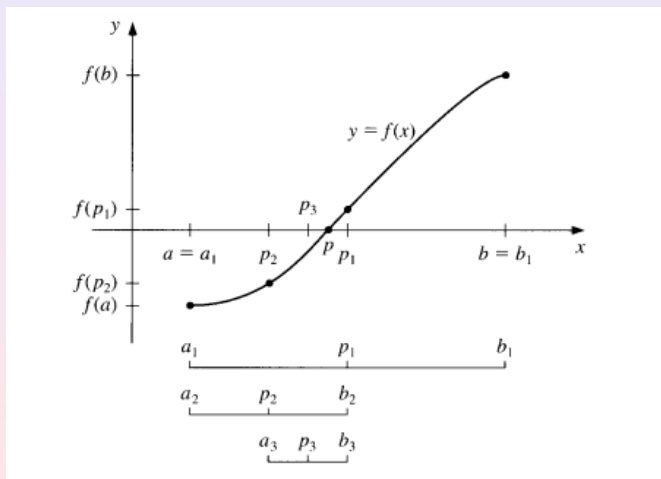
- 1 Bisection Method
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- 3 Newton's method
- 4 Error analysis for iterative methods
- 5 Accelerating convergence
- 6 Zeros of polynomials and Müller's method (SKIP)



# Bisection Method

## Idea

If  $f(x) \in C[a, b]$  and  $f(a)f(b) < 0$ , then  $\exists c \in (a, b)$  such that  $f(c) = 0$ .



## Bisection method algorithm

Given  $f(x)$  defined on  $(a, b)$ , the maximal number of iterations  $M$ , and stop criteria  $\delta$  and  $\varepsilon$ , this algorithm tries to locate one root of  $f(x)$ .

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Compute  $fa = f(a)$ ,  $fb = f(b)$ , and  $e = b - a$   
If  $sign(fa) = sign(fb)$ , then stop End If  
For  $k = 1, 2, \dots, M$   
     $e = e/2$ ,  $c = (a + b)/2$ ,  $fc = f(c)$   
    If  $(|e| < \delta$  or  $|fc| < \varepsilon)$ , then stop End If  
    If  $sign(fc) \neq sign(fa)$   
         $b = c$ ,  $fb = fc$   
    Else  
         $a = c$ ,  $fa = fc$   
    End If  
End For
```



Let  $\{c_n\}$  be the sequence of numbers produced. The algorithm should stop if one of the following conditions is satisfied.

- 1 the iteration number  $k > M$ ,
- 2  $|c_k - c_{k-1}| < \delta$ , or
- 3  $|f(c_k)| < \varepsilon$ .

Let  $[a_0, b_0], [a_1, b_1], \dots$  denote the successive intervals produced by the bisection algorithm. Then

$$\begin{aligned} & a = a_0 \leq a_1 \leq a_2 \leq \dots \leq b_0 = b \\ \Rightarrow & \{a_n\} \text{ and } \{b_n\} \text{ are bounded} \\ \Rightarrow & \lim_{n \rightarrow \infty} a_n \text{ and } \lim_{n \rightarrow \infty} b_n \text{ exist} \end{aligned}$$



Since

$$\begin{aligned}b_1 - a_1 &= \frac{1}{2}(b_0 - a_0) \\b_2 - a_2 &= \frac{1}{2}(b_1 - a_1) = \frac{1}{4}(b_0 - a_0) \\&\vdots \\b_n - a_n &= \frac{1}{2^n}(b_0 - a_0)\end{aligned}$$

hence

$$\lim_{n \rightarrow \infty} b_n - \lim_{n \rightarrow \infty} a_n = \lim_{n \rightarrow \infty} (b_n - a_n) = \lim_{n \rightarrow \infty} \frac{1}{2^n}(b_0 - a_0) = 0.$$

Therefore

$$\lim_{n \rightarrow \infty} a_n = \lim_{n \rightarrow \infty} b_n \equiv z.$$

Since  $f$  is a continuous function, we have that

$$\lim_{n \rightarrow \infty} f(a_n) = f\left(\lim_{n \rightarrow \infty} a_n\right) = f(z) \quad \text{and} \quad \lim_{n \rightarrow \infty} f(b_n) = f\left(\lim_{n \rightarrow \infty} b_n\right) = f(z).$$



On the other hand,

$$\begin{aligned} f(a_n)f(b_n) &< 0 \\ \Rightarrow \lim_{n \rightarrow \infty} f(a_n)f(b_n) &= f^2(z) \leq 0 \\ \Rightarrow f(z) &= 0 \end{aligned}$$

Therefore, the limit of the sequences  $\{a_n\}$  and  $\{b_n\}$  is a zero of  $f$  in  $[a, b]$ . Let  $c_n = \frac{1}{2}(a_n + b_n)$ . Then

$$\begin{aligned} |z - c_n| &= \left| \lim_{n \rightarrow \infty} a_n - \frac{1}{2}(a_n + b_n) \right| \\ &= \left| \frac{1}{2} \left[ \lim_{n \rightarrow \infty} a_n - b_n \right] + \frac{1}{2} \left[ \lim_{n \rightarrow \infty} a_n - a_n \right] \right| \\ &\leq \max \left\{ \left| \lim_{n \rightarrow \infty} a_n - b_n \right|, \left| \lim_{n \rightarrow \infty} a_n - a_n \right| \right\} \\ &\leq |b_n - a_n| = \frac{1}{2^n} |b_0 - a_0|. \end{aligned}$$

This proves the following theorem.



## Theorem

Let  $\{[a_n, b_n]\}$  denote the intervals produced by the bisection algorithm. Then  $\lim_{n \rightarrow \infty} a_n$  and  $\lim_{n \rightarrow \infty} b_n$  exist, are equal, and represent a zero of  $f(x)$ . If

$$z = \lim_{n \rightarrow \infty} a_n = \lim_{n \rightarrow \infty} b_n \quad \text{and} \quad c_n = \frac{1}{2}(a_n + b_n),$$

then

$$|z - c_n| \leq \frac{1}{2^n} (b_0 - a_0).$$

## Remark

$\{c_n\}$  converges to  $z$  with the rate of  $O(2^{-n})$ .





## Example

How many steps should be taken to compute a root of  $f(x) = x^3 + 4x^2 - 10 = 0$  on  $[1, 2]$  with relative error  $10^{-3}$ ?

*solution:* Seek an  $n$  such that

$$\frac{|z - c_n|}{|z|} \leq 10^{-3} \Rightarrow |z - c_n| \leq |z| \times 10^{-3}.$$

Since  $z \in [1, 2]$ , it is sufficient to show

$$|z - c_n| \leq 10^{-3}.$$

That is, we solve

$$2^{-n}(2 - 1) \leq 10^{-3} \Rightarrow -n \log_{10} 2 \leq -3$$

which gives  $n \geq 10$ .



# Fixed-Point Iteration

## Definition

$x$  is called a **fixed point** of a given function  $f$  if  $f(x) = x$ .

## Root-finding problems and fixed-point problems

- Find  $x^*$  such that  $f(x^*) = 0$ .

Let  $g(x) = x - f(x)$ . Then  $g(x^*) = x^* - f(x^*) = x^*$ .

$\Rightarrow x^*$  is a fixed point for  $g(x)$ .

- Find  $x^*$  such that  $g(x^*) = x^*$ .

Define  $f(x) = x - g(x)$  so that  $f(x^*) = x^* - g(x^*) = x^* - x^* = 0$

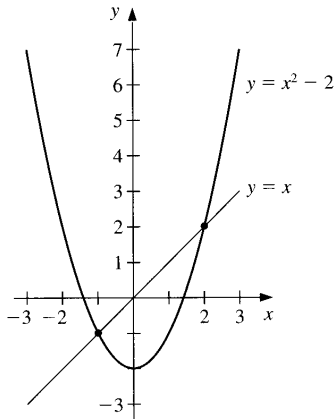
$\Rightarrow x^*$  is a zero of  $f(x)$ .



## Example

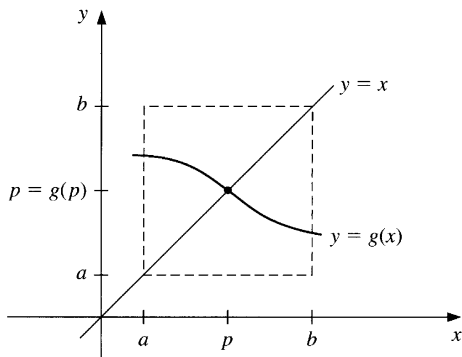
The function  $g(x) = x^2 - 2$ , for  $-2 \leq x \leq 3$ , has fixed points at  $x = -1$  and  $x = 2$  since

$$g(-1) = (-1)^2 - 2 = -1 \quad \text{and} \quad g(2) = 2^2 - 2 = 2.$$



## Theorem (Existence and uniqueness)

- 1 If  $g \in C[a, b]$  such that  $a \leq g(x) \leq b$  for all  $x \in [a, b]$ , then  $g$  has a fixed point in  $[a, b]$ .
- 2 If, in addition,  $g'(x)$  exists in  $(a, b)$  and there exists a positive constant  $M < 1$  such that  $|g'(x)| \leq M < 1$  for all  $x \in (a, b)$ . Then the fixed point is *unique*.



# Proof

*Existence:*

- If  $g(a) = a$  or  $g(b) = b$ , then  $a$  or  $b$  is a fixed point of  $g$  and we are done.
- Otherwise, it must be  $g(a) > a$  and  $g(b) < b$ . The function  $h(x) = g(x) - x$  is continuous on  $[a, b]$ , with

$$h(a) = g(a) - a > 0 \quad \text{and} \quad h(b) = g(b) - b < 0.$$

By the Intermediate Value Theorem,  $\exists x^* \in [a, b]$  such that  $h(x^*) = 0$ . That is

$$g(x^*) - x^* = 0 \Rightarrow g(x^*) = x^*.$$

Hence  $g$  has a fixed point  $x^*$  in  $[a, b]$ .



# Proof

*Uniqueness:*

Suppose that  $p \neq q$  are both fixed points of  $g$  in  $[a, b]$ . By the Mean-Value theorem, there exists  $\xi$  between  $p$  and  $q$  such that

$$g'(\xi) = \frac{g(p) - g(q)}{p - q} = \frac{p - q}{p - q} = 1.$$

However, this contradicts to the assumption that  $|g'(x)| \leq M < 1$  for all  $x$  in  $[a, b]$ . Therefore the fixed point of  $g$  is unique.  $\square$



## Example

Show that the following function has a unique fixed point.

$$g(x) = (x^2 - 1)/3, \quad x \in [-1, 1].$$

*Solution:* The Extreme Value Theorem implies that

$$\begin{aligned} \min_{x \in [-1, 1]} g(x) &= g(0) = -\frac{1}{3}, \\ \max_{x \in [-1, 1]} g(x) &= g(\pm 1) = 0. \end{aligned}$$

That is  $g(x) \in [-1/3, 0] \subset [-1, 1], \forall x \in [-1, 1]$ .

Moreover,  $g$  is continuous and

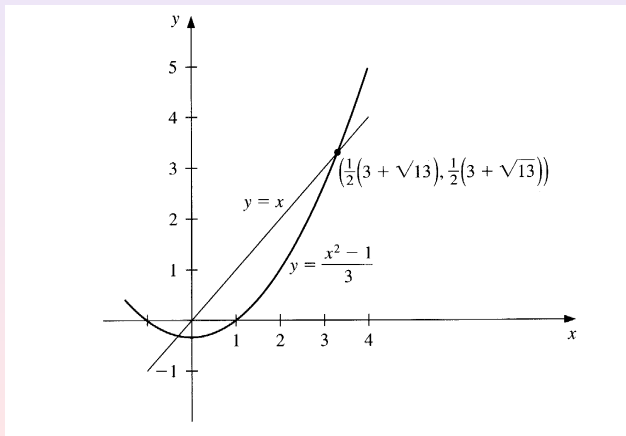
$$|g'(x)| = \left| \frac{2x}{3} \right| \leq \frac{2}{3}, \quad \forall x \in (-1, 1).$$

By above theorem,  $g$  has a unique fixed point in  $[-1, 1]$ .



Let  $p$  be such unique fixed point of  $g$ . Then

$$\begin{aligned} p = g(p) = \frac{p^2 - 1}{3} &\Rightarrow p^2 - 3p - 1 = 0 \\ &\Rightarrow p = \frac{1}{2}(3 + \sqrt{13}). \end{aligned}$$





## Fixed-point iteration or functional iteration

Given a continuous function  $g$ , choose an initial point  $x_0$  and generate  $\{x_k\}_{k=0}^{\infty}$  by

$$x_{k+1} = g(x_k), \quad k \geq 0.$$

$\{x_k\}$  may not converge, e.g.,  $g(x) = 3x$ . However, when the sequence converges, say,

$$\lim_{k \rightarrow \infty} x_k = x^*,$$

then, since  $g$  is continuous,

$$g(x^*) = g\left(\lim_{k \rightarrow \infty} x_k\right) = \lim_{k \rightarrow \infty} g(x_k) = \lim_{k \rightarrow \infty} x_{k+1} = x^*.$$

That is,  $x^*$  is a fixed point of  $g$ .



## Fixed-point iteration

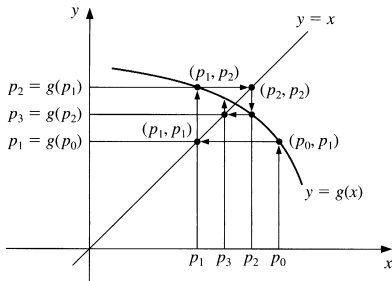
Given  $x_0$ , tolerance  $TOL$ , maximum number of iteration  $M$ .

Set  $i = 1$  and  $x = g(x_0)$ .

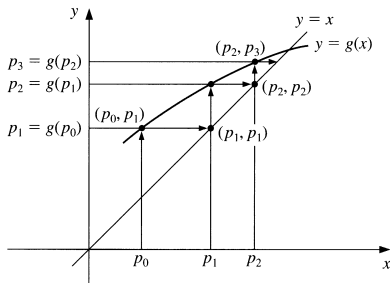
While  $i \leq M$  and  $|x - x_0| \geq TOL$

Set  $i = i + 1$ ,  $x_0 = x$  and  $x = g(x_0)$ .

End While



(a)



(b)

## Example

The equation

$$x^3 + 4x^2 - 10 = 0$$

has a unique root in  $[1, 2]$ . Change the equation to the fixed-point form  $x = g(x)$ .

$$(a) \quad x = g_1(x) \equiv x - f(x) = x - x^3 - 4x^2 + 10$$

$$(b) \quad x = g_2(x) = \left(\frac{10}{x} - 4x\right)^{1/2}$$

$$x^3 = 10 - 4x^2 \Rightarrow x^2 = \frac{10}{x} - 4x \Rightarrow x = \pm \left(\frac{10}{x} - 4x\right)^{1/2}$$



$$(c) x = g_3(x) = \frac{1}{2} (10 - x^3)^{1/2}$$

$$4x^2 = 10 - x^3 \Rightarrow x = \pm \frac{1}{2} (10 - x^3)^{1/2}$$

$$(d) x = g_4(x) = \left( \frac{10}{4+x} \right)^{1/2}$$

$$x^2(x+4) = 10 \Rightarrow x = \pm \left( \frac{10}{4+x} \right)^{1/2}$$

$$(e) x = g_5(x) = x - \frac{x^3 + 4x^2 - 10}{3x^2 + 8x}$$

$$x = g_5(x) \equiv x - \frac{f(x)}{f'(x)}$$



## Results of the fixed-point iteration with initial point $x_0 = 1.5$

$n$	(a)	(b)	(c)	(d)	(e)
0	1.5	1.5	1.5	1.5	1.5
1	-0.875	0.8165	1.286953768	1.348399725	1.373333333
2	6.732	2.9969	1.402540804	1.367376372	1.365262015
3	-469.7	$(-8.65)^{1/2}$	1.345458374	1.364957015	1.365230014
4	$1.03 \times 10^8$		1.375170253	1.365264748	1.365230013
5			1.360094193	1.365225594	
6			1.367846968	1.365230576	
7			1.363887004	1.365229942	
8			1.365916734	1.365230022	
9			1.364878217	1.365230012	
10			1.365410062	1.365230014	
15			1.365223680	1.365230013	
20			1.365230236		
25			1.365230006		
30			1.365230013		



## Theorem (Fixed-point Theorem)

Let  $g \in C[a, b]$  be such that  $g(x) \in [a, b]$  for all  $x \in [a, b]$ . Suppose that  $g'$  exists on  $(a, b)$  and that  $\exists k$  with  $0 < k < 1$  such that

$$|g'(x)| \leq k, \quad \forall x \in (a, b).$$

Then, for any number  $x_0$  in  $[a, b]$ ,

$$x_n = g(x_{n-1}), \quad n \geq 1,$$

converges to the unique fixed point  $x$  in  $[a, b]$ .



*Proof:* By the assumptions, a unique fixed point exists in  $[a, b]$ . Since  $g([a, b]) \subseteq [a, b]$ ,  $\{x_n\}_{n=0}^{\infty}$  is defined and  $x_n \in [a, b]$  for all  $n \geq 0$ . Using the Mean Values Theorem and the fact that  $|g'(x)| \leq k$ , we have

$$|x - x_n| = |g(x_{n-1}) - g(x)| = |g'(\xi_n)| |x - x_{n-1}| \leq k|x - x_{n-1}|,$$

where  $\xi_n \in (a, b)$ . It follows that

$$|x_n - x| \leq k|x_{n-1} - x| \leq k^2|x_{n-2} - x| \leq \cdots \leq k^n|x_0 - x|. \quad (1)$$

Since  $0 < k < 1$ , we have

$$\lim_{n \rightarrow \infty} k^n = 0$$

and

$$\lim_{n \rightarrow \infty} |x_n - x| \leq \lim_{n \rightarrow \infty} k^n |x_0 - x| = 0.$$

Hence,  $\{x_n\}_{n=0}^{\infty}$  converges to  $x$ .



## Corollary

If  $g$  satisfies the hypotheses of above theorem, then

$$|x - x_n| \leq k^n \max\{x_0 - a, b - x_0\}$$

and

$$|x_n - x| \leq \frac{k^n}{1 - k} |x_1 - x_0|, \quad \forall n \geq 1.$$

*Proof:* From (1),

$$|x_n - x| \leq k^n |x_0 - x| \leq k^n \max\{x_0 - a, b - x_0\}.$$

For  $n \geq 1$ , using the Mean Values Theorem,

$$|x_{n+1} - x_n| = |g(x_n) - g(x_{n-1})| \leq k |x_n - x_{n-1}| \leq \cdots \leq k^n |x_1 - x_0|.$$





Thus, for  $m > n \geq 1$ ,

$$\begin{aligned} |x_m - x_n| &= |x_m - x_{m-1} + x_{m-1} - \cdots + x_{n+1} - x_n| \\ &\leq |x_m - x_{m-1}| + |x_{m-1} - x_{m-2}| + \cdots + |x_{n+1} - x_n| \\ &\leq k^{m-1}|x_1 - x_0| + k^{m-2}|x_1 - x_0| + \cdots + k^n|x_1 - x_0| \\ &= k^n|x_1 - x_0| (1 + k + k^2 + \cdots + k^{m-n-1}). \end{aligned}$$

It implies that

$$\begin{aligned} |x - x_n| &= \lim_{m \rightarrow \infty} |x_m - x_n| \leq \lim_{m \rightarrow \infty} k^n |x_1 - x_0| \sum_{j=0}^{m-n-1} k^j \\ &\leq k^n |x_1 - x_0| \sum_{j=0}^{\infty} k^j = \frac{k^n}{1-k} |x_1 - x_0|. \end{aligned}$$



## Example

For previous example,  $f(x) = x^3 + 4x^2 - 10 = 0$ .

Let  $g_1(x) = x - x^3 - 4x^2 + 10$ , we have

$$g_1(1) = 6 \quad \text{and} \quad g_1(2) = -12,$$

so  $g_1([1, 2]) \not\subseteq [1, 2]$ . Moreover,

$$g_1'(x) = 1 - 3x^2 - 8x \quad \Rightarrow \quad |g_1'(x)| \geq 1 \quad \forall x \in [1, 2]$$

Convergence is NOT guaranteed. In fact, it almost for sure will not converge since when  $x_n$  is close to the solution  $x^*$ ,

$$|x_n - x^*| = |g_1(x_{n-1}) - g_1(x^*)| = |g_1'(c)(x_{n-1} - x^*)| > |x_{n-1} - x^*|.$$

The error is amplified whenever  $x_n$  is close to convergence. The only possibility for convergence is when  $x_n$  is far from  $x^*$  and (by chance, and very unlikely) that  $g_1(x_n) = x^*$ .

For  $g_3(x) = \frac{1}{2}(10 - x^3)^{1/2}$ ,  $\forall x \in [1, 1.5]$ ,

$$g_3'(x) = -\frac{3}{4}x^2(10 - x^3)^{-1/2} < 0, \forall x \in [1, 1.5],$$

so  $g_3$  is strictly decreasing on  $[1, 1.5]$  and

$$1 < 1.28 \approx g_3(1.5) \leq g_3(x) \leq g_3(1) = 1.5, \forall x \in [1, 1.5].$$

On the other hand,

$$|g_3'(x)| \leq |g_3'(1.5)| \approx 0.66, \forall x \in [1, 1.5]$$

Hence, the sequence is convergent to the fixed point.



For  $g_4(x) = \sqrt{10/(4+x)}$ , we have

$$\sqrt{\frac{10}{6}} \leq g_4(x) \leq \sqrt{\frac{10}{5}}, \quad \forall x \in [1, 2] \quad \Rightarrow \quad g_4([1, 2]) \subseteq [1, 2]$$

Moreover,

$$|g_4'(x)| = \left| \frac{-5}{\sqrt{10}(4+x)^{3/2}} \right| \leq \frac{5}{\sqrt{10}(5)^{3/2}} < 0.15, \quad \forall x \in [1, 2].$$

The bound of  $|g_4'(x)|$  is much smaller than the bound of  $|g_3'(x)|$ , which explains the more rapid convergence using  $g_4$ .



## Newton's method

Suppose that  $f : \mathbb{R} \rightarrow \mathbb{R}$  and  $f \in C^2[a, b]$ , i.e.,  $f''$  exists and is continuous. If  $f(x^*) = 0$  and  $x^* = x + h$  where  $h$  is small, then by Taylor's theorem

$$\begin{aligned}0 = f(x^*) &= f(x + h) \\ &= f(x) + f'(x)h + \frac{1}{2}f''(x)h^2 + \frac{1}{3!}f'''(x)h^3 + \dots \\ &= f(x) + f'(x)h + O(h^2).\end{aligned}$$

Since  $h$  is small,  $O(h^2)$  is negligible. It is reasonable to drop  $O(h^2)$  terms. This implies

$$f(x) + f'(x)h \approx 0 \quad \text{and} \quad h \approx -\frac{f(x)}{f'(x)}, \quad \text{if } f'(x) \neq 0.$$

Hence

$$x + h = x - \frac{f(x)}{f'(x)}$$

is a better approximation to  $x^*$ .



This sets the stage for the **Newton-Raphson's** method, which starts with an initial approximation  $x_0$  and generates the sequence  $\{x_n\}_{n=0}^{\infty}$  defined by

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}.$$

Since the Taylor's expansion of  $f(x)$  at  $x_k$  is given by

$$f(x) = f(x_k) + f'(x_k)(x - x_k) + \frac{1}{2}f''(x_k)(x - x_k)^2 + \cdots .$$

At  $x_k$ , one uses the **tangent line**

$$y = \ell(x) = f(x_k) + f'(x_k)(x - x_k)$$

to **approximate the curve** of  $f(x)$  and uses the zero of the tangent line to approximate the zero of  $f(x)$ .



## Newton's Method

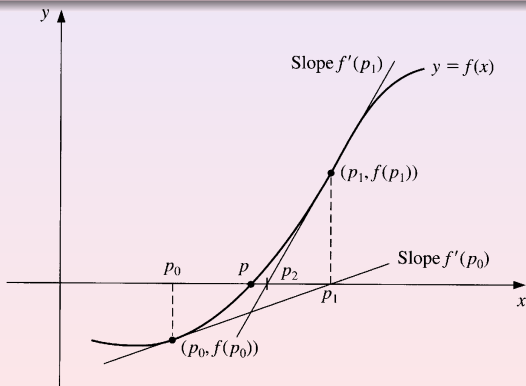
Given  $x_0$ , tolerance  $TOL$ , maximum number of iteration  $M$ .

Set  $i = 1$  and  $x = x_0 - f(x_0)/f'(x_0)$ .

While  $i \leq M$  and  $|x - x_0| \geq TOL$

Set  $i = i + 1$ ,  $x_0 = x$  and  $x = x_0 - f(x_0)/f'(x_0)$ .

End While



## Three stopping-technique inequalities

$$(a). \quad |x_n - x_{n-1}| < \varepsilon,$$

$$(b). \quad \frac{|x_n - x_{n-1}|}{|x_n|} < \varepsilon, \quad x_n \neq 0,$$

$$(c). \quad |f(x_n)| < \varepsilon.$$

Note that Newton's method for solving  $f(x) = 0$

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}, \quad \text{for } n \geq 1$$

is just a special case of functional iteration in which

$$g(x) = x - \frac{f(x)}{f'(x)}.$$



## Example

The following table shows the convergence behavior of Newton's method applied to solving  $f(x) = x^2 - 1 = 0$ . Observe the quadratic convergence rate.

$n$	$x_n$	$ e_n  \equiv  1 - x_n $
0	2.0	1
1	1.25	0.25
2	1.025	2.5e-2
3	1.0003048780488	3.048780488e-4
4	1.0000000464611	4.64611e-8
5	1.0	0



## Theorem

Assume  $f(x^*) = 0$ ,  $f'(x^*) \neq 0$  and  $f(x)$ ,  $f'(x)$  and  $f''(x)$  are *continuous* on  $N_\varepsilon(x^*)$ . Then if  $x_0$  is chosen *sufficiently close* to  $x^*$ , then

$$\left\{ x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)} \right\} \rightarrow x^*.$$

*Proof:* Define

$$g(x) = x - \frac{f(x)}{f'(x)}.$$

Find an interval  $[x^* - \delta, x^* + \delta]$  such that

$$g([x^* - \delta, x^* + \delta]) \subseteq [x^* - \delta, x^* + \delta]$$

and

$$|g'(x)| \leq k < 1, \quad \forall x \in (x^* - \delta, x^* + \delta).$$



Since  $f'$  is continuous and  $f'(x^*) \neq 0$ , it implies that  $\exists \delta_1 > 0$  such that  $f'(x) \neq 0 \forall x \in [x^* - \delta_1, x^* + \delta_1] \subseteq [a, b]$ . Thus,  $g$  is defined and continuous on  $[x^* - \delta_1, x^* + \delta_1]$ . Also

$$g'(x) = 1 - \frac{f'(x)f'(x) - f(x)f''(x)}{[f'(x)]^2} = \frac{f(x)f''(x)}{[f'(x)]^2},$$

for  $x \in [x^* - \delta_1, x^* + \delta_1]$ . Since  $f''$  is continuous on  $[a, b]$ , we have  $g'$  is continuous on  $[x^* - \delta_1, x^* + \delta_1]$ .

By assumption  $f(x^*) = 0$ , so

$$g'(x^*) = \frac{f(x^*)f''(x^*)}{|f'(x^*)|^2} = 0.$$

Since  $g'$  is continuous on  $[x^* - \delta_1, x^* + \delta_1]$  and  $g'(x^*) = 0$ ,  $\exists \delta$  with  $0 < \delta < \delta_1$  and  $k \in (0, 1)$  such that

$$|g'(x)| \leq k, \quad \forall x \in [x^* - \delta, x^* + \delta].$$



Claim:  $g([x^* - \delta, x^* + \delta]) \subseteq [x^* - \delta, x^* + \delta]$ .

If  $x \in [x^* - \delta, x^* + \delta]$ , then, by the Mean Value Theorem,  $\exists \xi$  between  $x$  and  $x^*$  such that

$$|g(x) - g(x^*)| = |g'(\xi)||x - x^*|.$$

It implies that

$$\begin{aligned} |g(x) - x^*| &= |g(x) - g(x^*)| = |g'(\xi)||x - x^*| \\ &\leq k|x - x^*| < |x - x^*| < \delta. \end{aligned}$$

Hence,  $g([x^* - \delta, x^* + \delta]) \subseteq [x^* - \delta, x^* + \delta]$ .

By the Fixed-Point Theorem, the sequence  $\{x_n\}_{n=0}^{\infty}$  defined by

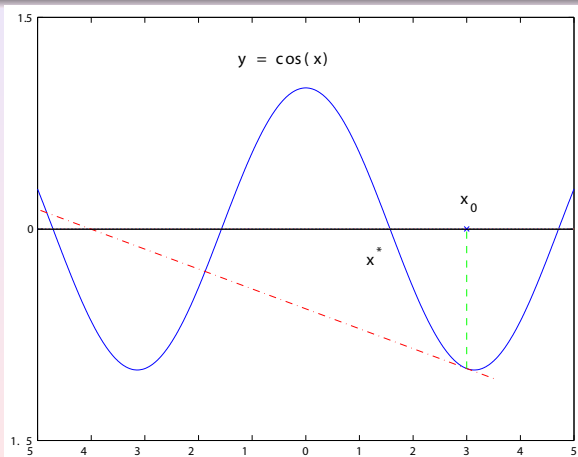
$$x_n = g(x_{n-1}) = x_{n-1} - \frac{f(x_{n-1})}{f'(x_{n-1})}, \quad \text{for } n \geq 1,$$

converges to  $x^*$  for any  $x_0 \in [x^* - \delta, x^* + \delta]$ .



## Example

When Newton's method applied to  $f(x) = \cos x$  with starting point  $x_0 = 3$ , which is close to the root  $\frac{\pi}{2}$  of  $f$ , it produces  $x_1 = -4.01525$ ,  $x_2 = -4.8526$ ,  $\dots$ , which converges to another root  $-\frac{3\pi}{2}$ .



# Secant method

## Disadvantage of Newton's method

In many applications, the derivative  $f'(x)$  is very expensive to compute, or the function  $f(x)$  is not given in an algebraic formula so that  $f'(x)$  is not available.

By definition,

$$f'(x_{n-1}) = \lim_{x \rightarrow x_{n-1}} \frac{f(x) - f(x_{n-1})}{x - x_{n-1}}.$$

Letting  $x = x_{n-2}$ , we have

$$f'(x_{n-1}) \approx \frac{f(x_{n-2}) - f(x_{n-1})}{x_{n-2} - x_{n-1}} = \frac{f(x_{n-1}) - f(x_{n-2})}{x_{n-1} - x_{n-2}}.$$

Using this approximation for  $f'(x_{n-1})$  in Newton's formula gives

$$x_n = x_{n-1} - \frac{f(x_{n-1})(x_{n-1} - x_{n-2})}{f(x_{n-1}) - f(x_{n-2})},$$

which is called the **Secant method**.



From geometric point of view, we use a **secant line** through  $x_{n-1}$  and  $x_{n-2}$  instead of the tangent line to approximate the function at the point  $x_{n-1}$ . The slope of the secant line is

$$s_{n-1} = \frac{f(x_{n-1}) - f(x_{n-2})}{x_{n-1} - x_{n-2}}$$

and the equation is

$$M(x) = f(x_{n-1}) + s_{n-1}(x - x_{n-1}).$$

The zero of the secant line

$$x = x_{n-1} - \frac{f(x_{n-1})}{s_{n-1}} = x_{n-1} - f(x_{n-1}) \frac{x_{n-1} - x_{n-2}}{f(x_{n-1}) - f(x_{n-2})}$$

is then used as a new approximate  $x_n$ .



## Secant Method

Given  $x_0, x_1$ , tolerance  $TOL$ , maximum number of iteration  $M$ .

Set  $i = 2$ ;  $y_0 = f(x_0)$ ;  $y_1 = f(x_1)$ ;

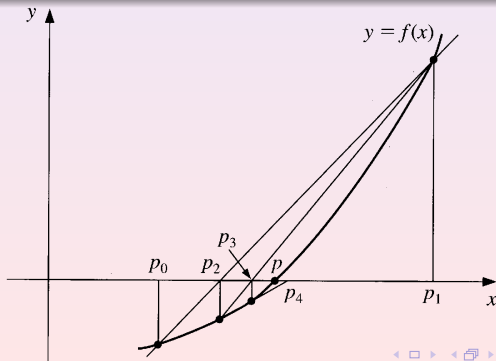
$$x = x_1 - y_1(x_1 - x_0)/(y_1 - y_0).$$

While  $i \leq M$  and  $|x - x_1| \geq TOL$

Set  $i = i + 1$ ;  $x_0 = x_1$ ;  $y_0 = y_1$ ;  $x_1 = x$ ;  $y_1 = f(x)$ ;

$$x = x_1 - y_1(x_1 - x_0)/(y_1 - y_0).$$

End While





## Method of False Position

- 1 Choose initial approximations  $x_0$  and  $x_1$  with  $f(x_0)f(x_1) < 0$ .
- 2  $x_2 = x_1 - f(x_1)(x_1 - x_0)/(f(x_1) - f(x_0))$
- 3 Decide which secant line to use to compute  $x_3$ :  
If  $f(x_2)f(x_1) < 0$ , then  $x_1$  and  $x_2$  bracket a root, i.e.,

$$x_3 = x_2 - f(x_2)(x_2 - x_1)/(f(x_2) - f(x_1))$$

Else,  $x_0$  and  $x_2$  bracket a root, i.e.,

$$x_3 = x_2 - f(x_2)(x_2 - x_0)/(f(x_2) - f(x_0))$$

End if



## Method of False Position

Given  $x_0, x_1$ , tolerance  $TOL$ , maximum number of iteration  $M$ .

Set  $i = 2$ ;  $y_0 = f(x_0)$ ;  $y_1 = f(x_1)$ ;  $x = x_1 - y_1(x_1 - x_0)/(y_1 - y_0)$ .

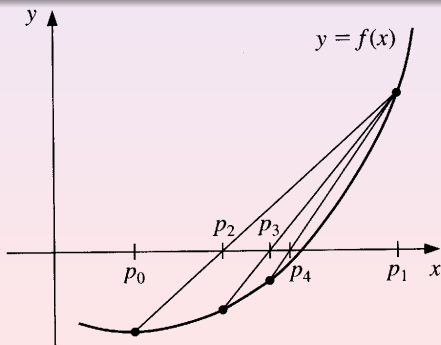
While  $i \leq M$  and  $|x - x_1| \geq TOL$

Set  $i = i + 1$ ;  $y = f(x)$ .

If  $y \cdot y_1 < 0$ , then set  $x_0 = x$ ;  $y_0 = y_1$ .

Set  $x_1 = x$ ;  $y_1 = y$ ;  $x = x_1 - y_1(x_1 - x_0)/(y_1 - y_0)$ .

End While



# Error analysis for iterative methods

## Definition

Let  $\{x_n\} \rightarrow x^*$ . If there are positive constants  $c$  and  $\alpha$  such that

$$\lim_{n \rightarrow \infty} \frac{|x_{n+1} - x^*|}{|x_n - x^*|^\alpha} = c,$$

then we say the **rate of convergence** is of **order  $\alpha$** .

We say that the rate of convergence is

① **linear** if  $\alpha = 1$  and  $0 < c < 1$ .

② **superlinear** if

$$\lim_{n \rightarrow \infty} \frac{|x_{n+1} - x^*|}{|x_n - x^*|} = 0;$$

③ **quadratic** if  $\alpha = 2$ .

Suppose that  $\{x_n\}_{n=0}^{\infty}$  and  $\{\tilde{x}_n\}_{n=0}^{\infty}$  are linearly and quadratically convergent to  $x^*$ , respectively, with the same constant  $c = 0.5$ . For simplicity, suppose that

$$\frac{|x_{n+1} - x^*|}{|x_n - x^*|} \approx c \quad \text{and} \quad \frac{|\tilde{x}_{n+1} - x^*|}{|\tilde{x}_n - x^*|^2} \approx c.$$

These imply that

$$|x_n - x^*| \approx c|x_{n-1} - x^*| \approx c^2|x_{n-2} - x^*| \approx \cdots \approx c^n|x_0 - x^*|,$$

and

$$\begin{aligned} |\tilde{x}_n - x^*| &\approx c|\tilde{x}_{n-1} - x^*|^2 \approx c [c|\tilde{x}_{n-2} - x^*|^2]^2 = c^3|\tilde{x}_{n-2} - x^*|^4 \\ &\approx c^3 [c|\tilde{x}_{n-3} - x^*|^2]^4 = c^7|\tilde{x}_{n-3} - x^*|^8 \\ &\approx \cdots \approx c^{2^n - 1}|\tilde{x}_0 - x^*|^{2^n}. \end{aligned}$$



## Remark

Quadratically convergent sequences generally converge much more quickly than those that converge only linearly.

## Theorem

Let  $g \in C[a, b]$  with  $g([a, b]) \subseteq [a, b]$ . Suppose that  $g'$  is continuous on  $(a, b)$  and  $\exists k \in (0, 1)$  such that

$$|g'(x)| \leq k, \quad \forall x \in (a, b).$$

If  $g'(x^*) \neq 0$ , then for any  $x_0 \in [a, b]$ , the sequence

$$x_n = g(x_{n-1}), \quad \text{for } n \geq 1$$

converges only linearly to the unique fixed point  $x^*$  in  $[a, b]$ .



*Proof:*

- By the Fixed-Point Theorem, the sequence  $\{x_n\}_{n=0}^{\infty}$  converges to  $x^*$ .
- Since  $g'$  exists on  $(a, b)$ , by the Mean Value Theorem,  $\exists \xi_n$  between  $x_n$  and  $x^*$  such that

$$x_{n+1} - x^* = g(x_n) - g(x^*) = g'(\xi_n)(x_n - x^*).$$

- $\because \{x_n\}_{n=0}^{\infty} \rightarrow x^* \Rightarrow \{\xi_n\}_{n=0}^{\infty} \rightarrow x^*$
- Since  $g'$  is continuous on  $(a, b)$ , we have

$$\lim_{n \rightarrow \infty} g'(\xi_n) = g'(x^*).$$

- Thus,

$$\lim_{n \rightarrow \infty} \frac{|x_{n+1} - x^*|}{|x_n - x^*|} = \lim_{n \rightarrow \infty} |g'(\xi_n)| = |g'(x^*)|.$$

Hence, if  $g'(x^*) \neq 0$ , fixed-point iteration exhibits linear convergence.



## Theorem

Let  $x^*$  be a fixed point of  $g$  and  $I$  be an open interval with  $x^* \in I$ . Suppose that  $g'(x^*) = 0$  and  $g''$  is continuous with

$$|g''(x)| < M, \quad \forall x \in I.$$

Then  $\exists \delta > 0$  such that

$$\{x_n = g(x_{n-1})\}_{n=1}^{\infty} \rightarrow x^* \quad \text{for } x_0 \in [x^* - \delta, x^* + \delta]$$

at least quadratically. Moreover,

$$|x_{n+1} - x^*| < \frac{M}{2} |x_n - x^*|^2, \quad \text{for sufficiently large } n.$$



*Proof:*

- Since  $g'(x^*) = 0$  and  $g'$  is continuous on  $I$ ,  $\exists \delta$  such that  $[x^* - \delta, x^* + \delta] \subset I$  and

$$|g'(x)| \leq k < 1, \quad \forall x \in [x^* - \delta, x^* + \delta].$$

- In the proof of the convergence for Newton's method, we have

$$\{x_n\}_{n=0}^{\infty} \subset [x^* - \delta, x^* + \delta].$$

- Consider the Taylor expansion of  $g(x_n)$  at  $x^*$

$$\begin{aligned}x_{n+1} = g(x_n) &= g(x^*) + g'(x^*)(x_n - x^*) + \frac{g''(\xi)}{2}(x_n - x^*)^2 \\ &= x^* + \frac{g''(\xi)}{2}(x_n - x^*)^2,\end{aligned}$$

where  $\xi$  lies between  $x_n$  and  $x^*$ .





- Since

$$|g'(x)| \leq k < 1, \quad \forall x \in [x^* - \delta, x^* + \delta]$$

and

$$g([x^* - \delta, x^* + \delta]) \subseteq [x^* - \delta, x^* + \delta],$$

it follows that  $\{x_n\}_{n=0}^{\infty}$  converges to  $x^*$ .

- But  $\xi_n$  is between  $x_n$  and  $x^*$  for each  $n$ , so  $\{\xi_n\}_{n=0}^{\infty}$  also converges to  $x^*$  and

$$\lim_{n \rightarrow \infty} \frac{|x_{n+1} - x^*|}{|x_n - x^*|^2} = \frac{|g''(x^*)|}{2} < \frac{M}{2}.$$

- It implies that  $\{x_n\}_{n=0}^{\infty}$  is quadratically convergent to  $x^*$  if  $g''(x^*) \neq 0$  and

$$|x_{n+1} - x^*| < \frac{M}{2} |x_n - x^*|^2, \quad \text{for sufficiently large } n. \quad \square$$



## Example

Recall that Newton's method  $x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}$  corresponds to  $g(x) = x - \frac{f(x)}{f'(x)}$ . Suppose that  $f(x)$  has a  $m$ -fold root at  $x^*$ , that is

$$f(x) = (x - x^*)^m q(x), \quad q(x^*) \neq 0.$$

Let  $\mu(x) = \frac{f(x)}{f'(x)} = (x - x^*) \frac{q(x)}{mq(x) + (x - x^*)q'(x)}$ , it is easy to see that  $\mu'(x^*) = \frac{1}{m}$ . It follows that  $0 \leq g'(x_*) = 1 - \frac{1}{m} < 1$ . Hence Newton's method is locally convergent. Moreover, it converges **quadratically** for simple roots ( $m = 1$ ) and **linearly** for multiple roots ( $m > 1$ ).

*Remedy for slow convergence on multiple roots ( $m > 1$ ):*

- If  $m$  is known, take  $x_{n+1} = x_n - \frac{mf(x_n)}{f'(x_n)}$ .
- If  $m$  is not known, take  $x_{n+1} = x_n - \frac{\mu(x_n)}{\mu'(x_n)}$ , since  $\mu(x) = \frac{f(x)}{f'(x)} = \frac{O(x-x^*)^m}{O(x-x^*)^{m-1}} = O(x-x^*)$  always has a simple root at  $x^*$  for any  $m \geq 1$ . This is known as modified Newton's method.



# Global Convergence for Convex (Concave) Functions

## Theorem

If  $f \in C^2$ ,  $f'' > 0$  and  $f(x) = 0$  has a root, then Newton's method always converges to a root  $x^*$  for any initial  $x_0$ .

*Proof:*

It suffices to consider the case where  $f' > 0$ ,  $f'' > 0$  and  $f(x) = 0$  has a root. In this case, the root  $x^*$  is unique. Define  $e_n = x_n - x^*$ . Since  $x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}$ . It follows that

$$e_{n+1} = e_n - \frac{f(x_n)}{f'(x_n)}. \quad (2)$$

Moreover, since  $f(x^*) = f(x_n) + f'(x_n)(x^* - x_n) + \frac{f''(\xi_n)}{2}(x^* - x_n)^2$ , we also have  $f(x_n) = f'(x_n)e_n - \frac{f''(\xi_n)}{2}e_n^2$ . Therefore

$$e_{n+1} = e_n - \frac{f(x_n)}{f'(x_n)} = \frac{f''(\xi_n)}{2f'(x_n)}e_n^2 > 0. \quad (3)$$

Consequently  $x_{n+1} > x^*$  and  $f(x_{n+1}) > 0$  for all  $n \geq 0$ .



Moreover  $e_{n+1} = e_n - \frac{f(x_n)}{f'(x_n)} < e_n$ , we conclude that

$$0 < \dots < x_{n+1} < x_n < \dots < x_1$$

and  $x_n$  converges monotonically to some  $\tilde{x}$  satisfying  $\tilde{x} = \tilde{x} - \frac{f(\tilde{x})}{f'(\tilde{x})}$ , that is  $f(\tilde{x}) = 0$ , thus  $\tilde{x} = x^*$  by uniqueness of the root.

The proof for other cases

- $f' < 0$ ,  $f'' > 0$ ,  $f(x) = 0$  has a root.
- $f'' > 0$ , has two distinct roots.
- $f'' > 0$ , has a double root.

are similar. So is the concave case ( $f'' < 0$ ).



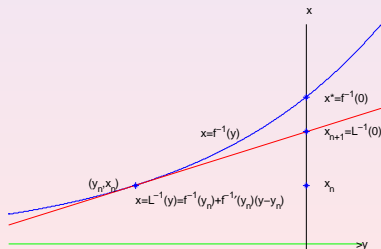
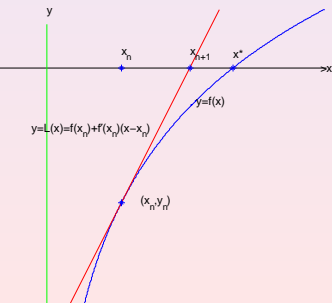
## Alternative Error Estimate for Newton's Method

Suppose  $f'(x^*) \neq 0$ , then both  $f(x)$  and its linearization at  $(x_n, y_n)$ ,  $L_n(x)$ , are locally invertible (Inverse Function Theorem). The formula of the tangent lines are given by

$$L_n(x) = f(x_n) + \frac{df(x_n)}{dx}(x - x_n)$$

and

$$L_n^{-1}(y) = f^{-1}(y_n) + \frac{df^{-1}(y_n)}{dy}(y - y_n) = x_n + \frac{1}{f'(x_n)}(y - y_n)$$



Since  $x^* = f^{-1}(0)$  and  $x_{n+1} = L_n^{-1}(0)$ , the error estimate for Newton's method reduces to error estimate between  $f^{-1}(y)$  and its linearization approximation  $L_n^{-1}(y)$  at  $y = 0$ . From standard analysis, the error is proportional to  $(0 - y_n)^2$ :

$$\begin{aligned} |x_{n+1} - x^*| &= |L_n^{-1}(0) - f^{-1}(0)| = \frac{1}{2} \left| \frac{d^2 f^{-1}}{dy^2}(\eta_n)(y_n - 0)^2 \right| \\ &= \frac{1}{2} \left| \frac{d^2 f^{-1}}{dy^2}(\eta_n) \right| (f(x_n) - f(x^*))^2 = \left( \frac{1}{2} \left| \frac{d^2 f^{-1}}{dy^2}(\eta_n) \right| \cdot (f'(\xi_n))^2 \right) (x_n - x^*)^2 \end{aligned}$$

The main advantage of this formulation:

Higher order approximations of  $f^{-1}(0)$ , such as quadratic approximation, gives rise to higher order iteration schemes for solving the original equation  $f(x) = 0$ .



# Error Analysis of Secant Method

*Reference:* D. Kincaid and W. Cheney, "Numerical analysis"

Let  $x^*$  denote the exact solution of  $f(x) = 0$ ,  $e_k = x_k - x^*$  be the errors at the  $k$ -th step. Then

$$\begin{aligned}e_{k+1} &= x_{k+1} - x^* \\&= x_k - f(x_k) \frac{x_k - x_{k-1}}{f(x_k) - f(x_{k-1})} - x^* \\&= \frac{1}{f(x_k) - f(x_{k-1})} [(x_{k-1} - x^*)f(x_k) - (x_k - x^*)f(x_{k-1})] \\&= \frac{1}{f(x_k) - f(x_{k-1})} (e_{k-1}f(x_k) - e_k f(x_{k-1})) \\&= e_k e_{k-1} \left( \frac{\frac{1}{e_k} f(x_k) - \frac{1}{e_{k-1}} f(x_{k-1})}{x_k - x_{k-1}} \cdot \frac{x_k - x_{k-1}}{f(x_k) - f(x_{k-1})} \right)\end{aligned}$$



To estimate the numerator  $\frac{1}{e_k} f(x_k) - \frac{1}{e_{k-1}} f(x_{k-1})$ , we apply the Taylor's theorem

$$f(x_k) = f(x^* + e_k) = f(x^*) + f'(x^*)e_k + \frac{1}{2}f''(x^*)e_k^2 + O(e_k^3),$$

to get

$$\frac{1}{e_k} f(x_k) = f'(x^*) + \frac{1}{2}f''(x^*)e_k + O(e_k^2).$$

Similarly,

$$\frac{1}{e_{k-1}} f(x_{k-1}) = f'(x^*) + \frac{1}{2}f''(x^*)e_{k-1} + O(e_{k-1}^2).$$

Hence

$$\frac{1}{e_k} f(x_k) - \frac{1}{e_{k-1}} f(x_{k-1}) \approx \frac{1}{2}(e_k - e_{k-1})f''(x^*).$$

Since  $x_k - x_{k-1} = e_k - e_{k-1}$  and

$$\frac{x_k - x_{k-1}}{f(x_k) - f(x_{k-1})} \rightarrow \frac{1}{f'(x^*)},$$





we have

$$\begin{aligned} e_{k+1} &\approx e_k e_{k-1} \left( \frac{\frac{1}{2}(e_k - e_{k-1})f''(x^*)}{e_k - e_{k-1}} \cdot \frac{1}{f'(x^*)} \right) = \frac{1}{2} \frac{f''(x^*)}{f'(x^*)} e_k e_{k-1} \\ &\equiv C e_k e_{k-1} \end{aligned} \quad (4)$$

To estimate the convergence rate, we assume

$$|e_{k+1}| \approx \eta |e_k|^\alpha,$$

where  $\eta > 0$  and  $\alpha > 0$  are constants, i.e.,

$$\frac{|e_{k+1}|}{\eta |e_k|^\alpha} \rightarrow 1 \quad \text{as } k \rightarrow \infty.$$

Then  $|e_k| \approx \eta |e_{k-1}|^\alpha$  which implies  $|e_{k-1}| \approx \eta^{-1/\alpha} |e_k|^{1/\alpha}$ . Hence (4) gives

$$\eta |e_k|^\alpha \approx C |e_k| \eta^{-1/\alpha} |e_k|^{1/\alpha} \implies C^{-1} \eta^{1+\frac{1}{\alpha}} \approx |e_k|^{1-\alpha+\frac{1}{\alpha}}.$$

Since  $|e_k| \rightarrow 0$  as  $k \rightarrow \infty$ , and  $C^{-1} \eta^{1+\frac{1}{\alpha}}$  is a nonzero constant,

$$1 - \alpha + \frac{1}{\alpha} = 0 \implies \alpha = \frac{1 + \sqrt{5}}{2} \approx 1.62.$$



This result implies that  $C^{-1}\eta^{1+\frac{1}{\alpha}} \rightarrow 1$  and

$$\eta \rightarrow C^{\frac{\alpha}{1+\alpha}} = \left( \frac{f''(x^*)}{2f'(x^*)} \right)^{0.62}.$$

In summary, we have shown that

$$|e_{k+1}| = \eta|e_k|^\alpha, \quad \alpha \approx 1.62,$$

that is, the **rate of convergence** is **superlinear**.

Rate of convergence:

- **secant** method: **superlinear**
- **Newton's** method: **quadratic**
- **bisection** method: **linear**



Each iteration of method requires

- secant method: one function evaluation
- Newton's method: two function evaluation, namely,  $f(x_k)$  and  $f'(x_k)$ .  
⇒ two steps of secant method are comparable to one step of Newton's method. Thus

$$|e_{k+2}| \approx \eta |e_{k+1}|^\alpha \approx \eta^{1+\alpha} |e_k|^{\frac{3+\sqrt{5}}{2}} \approx \eta^{1+\alpha} |e_k|^{2.62}.$$

⇒ secant method is more efficient than Newton's method.

### Remark

Two steps of secant method would require a little more work than one step of Newton's method.



# Accelerating convergence

## Aitken's $\Delta^2$ method

- Accelerate the convergence of a sequence that is **linearly convergent**.
- Suppose  $\{x_n\}_{n=0}^{\infty}$  is a linearly convergent sequence with limit  $y$ . Construct  $\{\hat{x}_n\}_{n=0}^{\infty}$  that converges more rapidly to  $x$  than  $\{x_n\}_{n=0}^{\infty}$ .

For  $n$  sufficiently large,

$$\frac{x_{n+1} - x}{x_n - x} \approx \frac{x_{n+2} - x}{x_{n+1} - x}.$$

Then

$$(x_{n+1} - x)^2 \approx (x_{n+2} - x)(x_n - x),$$

so

$$x_{n+1}^2 - 2x_{n+1}x + x^2 \approx x_{n+2}x_n - (x_{n+2} + x_n)x + x^2$$



and

$$(x_{n+2} + x_n - 2x_{n+1})x \approx x_{n+2}x_n - x_{n+1}^2.$$

Solving for  $x$  gives

$$\begin{aligned}x &\approx \frac{x_{n+2}x_n - x_{n+1}^2}{x_{n+2} - 2x_{n+1} + x_n} \\&= \frac{x_n x_{n+2} - 2x_n x_{n+1} + x_n^2 - x_n^2 + 2x_n x_{n+1} - x_{n+1}^2}{x_{n+2} - 2x_{n+1} + x_n} \\&= \frac{x_n(x_{n+2} - 2x_{n+1} + x_n) - (x_{n+1} - x_n)^2}{(x_{n+2} - x_{n+1}) - (x_{n+1} - x_n)} \\&= x_n - \frac{(x_{n+1} - x_n)^2}{(x_{n+2} - x_{n+1}) - (x_{n+1} - x_n)}.\end{aligned}$$

Aitken's  $\Delta^2$  method

$$\hat{x}_n = x_n - \frac{(x_{n+1} - x_n)^2}{(x_{n+2} - x_{n+1}) - (x_{n+1} - x_n)} := \{\Delta^2\}x_n. \quad (5)$$

## Theorem

Suppose  $\{x_n\}_{n=0}^{\infty} \rightarrow x$  linearly and

$$\lim_{n \rightarrow \infty} \frac{x_{n+1} - x}{x_n - x} < 1.$$

Then  $\{\hat{x}_n\}_{n=0}^{\infty} \rightarrow x$  faster than  $\{x_n\}_{n=0}^{\infty}$  in the sense that

$$\lim_{n \rightarrow \infty} \frac{\hat{x}_n - x}{x_n - x} = 0.$$

- Aitken's  $\Delta^2$  method constructs the terms in order:

$$\begin{aligned} x_1 = g(x_0), \quad x_2 = g(x_1), \quad \hat{x}_0 = \{\Delta^2\}(x_0), \quad x_3 = g(x_2), \\ \hat{x}_1 = \{\Delta^2\}(x_1), \quad \hat{x}_2 = \{\Delta^2\}(x_2), \quad \dots, \end{aligned}$$

This is based on the assumption that  $|\hat{x}_0 - x| < |x_2 - x|$ ,  
 $|\hat{x}_1 - x| < |x_3 - x|$ , etc.



## Example

The sequence  $\{x_n = \cos(\frac{1}{n})\}_{n=1}^{\infty}$  converges linearly to  $x = 1$ .

$n$	$x_n$	$e_n$	$\hat{x}_n$	$\hat{e}_n$
1	0.54030	0.45969	0.96178	0.03822
2	0.87758	0.12242	0.98213	0.01787
3	0.94496	0.05504	0.98979	0.01021
4	0.96891	0.03109	0.99342	0.00658
5	0.98007	0.01993	0.99541	0.00459
6	0.98614	0.01386		
7	0.98981	0.01019		

- Note that  $\lim_{n \rightarrow \infty} \frac{x_{n+1} - x}{x_n - x} = 1$ . The assumption in previous Theorem is not satisfied. In this case,  $\{\hat{x}_n\}_{n=1}^{\infty}$  converges more rapidly to  $x = 1$  than  $\{x_{n+2}\}_{n=1}^{\infty}$ , but is of the same order. In fact  $\hat{e}_n/e_{n+2} \sim 1/3$  for large  $n$ . (Why?)



- Steffensen's method constructs the terms in order:

$$\begin{aligned}
 x_0^{(0)} &= x_0, & x_1^{(0)} (= x_1) &= g(x_0^{(0)}), & x_2^{(0)} (= x_2) &= g(x_1^{(0)}), \\
 x_0^{(1)} (= x_3) &= \{\Delta^2\}(x_0^{(0)}), & x_1^{(1)} (= x_4) &= g(x_0^{(1)}), & x_2^{(1)} (= x_5) &= g(x_1^{(1)}), \\
 x_0^{(2)} (= x_6) &= \{\Delta^2\}(x_0^{(1)}), & x_1^{(2)} (= x_7) &= g(x_0^{(2)}), & x_2^{(2)} (= x_8) &= g(x_1^{(2)}), \\
 &\dots,
 \end{aligned}$$

### Steffensen's method (To find a solution of $x = g(x)$ )

Given  $x_0$ , tolerance  $TOL$ , maximum number of iteration  $M$ .

Set  $i = 1$ .

While  $i \leq M$

Set  $x_1 = g(x_0)$ ;  $x_2 = g(x_1)$ ;  $x = x_0 - (x_1 - x_0)^2 / (x_2 - 2x_1 + x_0)$ .

If  $|x - x_0| < Tol$ , then STOP.

Set  $i = i + 1$ ;  $x_0 = x$ .

End While



## Theorem

Suppose that  $x = g(x)$  has the solution  $x^*$  with  $g'(x^*) \neq 1$ . If  $\exists \delta > 0$  such that  $g \in C^3[x^* - \delta, x^* + \delta]$ , then Steffensen's method gives quadratic convergence for any  $x_0 \in [x^* - \delta, x^* + \delta]$ .

*Proof:*

We denote by  $x_0, x_1, x_2, \dots$ , (instead of  $x_0^{(i)}, x_1^{(i)}, x_2^{(i)}$ ), the sequence generated by Steffensen's method.

We will show that  $|x_3 - x| \leq C|x_0 - x|^2$ ,  $|x_6 - x| \leq C|x_3 - x|^2$ , etc. to establish quadratic convergence. Denote by  $\Delta_i = x_i - x^*$ , we have

$$\begin{aligned}\Delta_1 &= x_1 - x^* = g(x_0) - g(x^*) \\ &= g'(x^*)(x_0 - x^*) + \frac{g''(x^*)}{2}(x_0 - x^*)^2 + O(\Delta_0^3) \\ \Delta_2 &= x_2 - x^* = g(x_1) - g(x^*) \\ &= g'(x^*)(x_1 - x^*) + \frac{g''(x^*)}{2}(x_1 - x^*)^2 + O(\Delta_1^3) \\ &= g'(x^*)^2 \Delta_0 + \left( \frac{g'(x^*)g''(x^*)}{2} + \frac{g''(x^*)g'(x^*)^2}{2} \right) \Delta_0^2 + O(\Delta_0^3)\end{aligned}$$



$$\begin{aligned}
x_3 &= x_0 - \frac{(x_1 - x_0)^2}{x_0 - 2x_1 + x_2} \\
\Delta_3 &= \Delta_0 - \frac{(\Delta_1 - \Delta_0)^2}{\Delta_0 - 2\Delta_1 + \Delta_2} \\
&= \Delta_0 - \frac{((g'(x) - 1)\Delta_0 + \frac{g''(x)}{2}\Delta_0^2 + O(\Delta_0^3))^2}{(g'^2(x) - 2g'(x) + 1)\Delta_0 + \frac{g''(x)}{2}(g'^2(x) + g'(x) - 2)\Delta_0^2 + O(\Delta_0^3)} \\
&= \Delta_0 - \Delta_0 \left( \frac{(g'(x) - 1)^2 + g''(x)(g'(x) - 1)\Delta_0 + O(\Delta_0^2)}{(g'(x) - 1)^2 + \frac{g''(x)}{2}(g'(x) + 2)(g'(x) - 1)\Delta_0 + O(\Delta_0^2)} \right) \\
&= \Delta_0 - \Delta_0 \left( 1 - \frac{g'(x^*)g''(x^*)}{2(g'(x^*) - 1)}\Delta_0 + O(\Delta_0^2) \right), \quad \text{if } g'(x^*) \neq 1
\end{aligned}$$

It follows that  $x_3 - x \approx C(x_0 - x)^2$ ,  $x_6 - x \approx C(x_3 - x)^2$ , etc. with  $C = \frac{g'(x^*)g''(x^*)}{2(g'(x^*) - 1)}$  if  $g'(x^*) \neq 1$ .



# Zeros of polynomials and Müller's method (SKIP)

- Horner's method:

Goal: Find successively all roots of a polynomial

$$P(x) = a_0 + a_1x + a_2x^2 + \cdots + a_{n-1}x^{n-1} + a_nx^n \quad (6)$$

with minimal computational cost.

The key step is to efficiently compute the quotient  $P(x)/(x - x^*)$  when a root  $x^*$  of  $P(x)$  has been found (eg. by Newton's method), or more generally, to find the quotient  $Q(x)$  and the remainder  $b_0$  such that

$$P(x) = (x - x_k)Q(x) + b_0, \quad (7)$$

for any given  $x_k$ . As a byproduct, one obtains  $P'(x_k) = Q(x_k)$  from (7) which can be utilized in the Newton-Raphson iteration  $x_{k+1} = x_k - \frac{P(x_k)}{P'(x_k)}$ . The coefficients of  $Q(x)$  can be obtained by assuming

$$Q(x) = b_1 + b_2x + \cdots + b_nx^{n-1}$$

and then comparing the coefficients in (6) and (7).



We have

$$\begin{aligned} b_0 + (x - x_k)Q(x) &= b_0 + (x - x_k)(b_1 + b_2x + \cdots + b_nx^{n-1}) \\ &= (b_0 - b_1x_k) + (b_1 - b_2x_k)x + \cdots + (b_{n-1} - b_nx_k)x^{n-1} + b_nx^n \\ &= a_0 + a_1x + \cdots + a_nx^n = P(x). \end{aligned}$$

and therefore

$$\begin{aligned} b_n &= a_n, \\ b_j &= a_j + b_{j+1}x_k, \quad \text{for } j = n-1, n-2, \dots, 1, 0, \end{aligned}$$

Moreover, the evaluation of  $Q(x_k)$  can be obtained through the nested expression:

$$Q(x) = b_1 + x(b_2 + x(b_3 + \cdots + x(b_{n-1} + xb_n)))$$

that is, let  $c_n = b_n (= a_n)$ , and for  $j = n-1, n-2, \dots, 1$ ,

$$c_j = b_j + c_{j+1}x_k,$$

then  $Q(x_k) = c_1$ .



## Horner's method (Evaluate $P(x_k)$ and $P'(x_k) = Q(x_k)$ )

Set  $y = a_n$ ;  $z = a_n$  ( $b_n = a_n$ ;  $c_n = a_n$ ).

For  $j = n - 1, n - 2, \dots, 1$

Set  $y = a_j + yx_k$ ;  $z = y + zx_k$  ( $b_j = a_j + b_{j+1}x_k$ ;  $c_j = b_j + c_{j+1}x_k$ ).

End for

Set  $y = a_0 + yx_k$  ( $b_0 = a_0 + b_1x_k$ ).

Output  $P(x_k) = y$  ( $= b_0$ );  $P'(x_k) = z$  ( $= c_1$ ).

If  $x_N$  is an approximate zero of  $P$ , then

$$\begin{aligned} P(x) &= (x - x_N)Q(x) + b_0 = (x - x_N)Q(x) + P(x_N) \\ &\approx (x - x_N)Q(x) \equiv (x - \hat{x}_1)Q_1(x). \end{aligned}$$

So  $x - \hat{x}_1$  is an approximate factor of  $P(x)$  and we can find a second approximate zero of  $P$  by applying Newton's method to  $Q_1(x)$ . The procedure is called deflation.

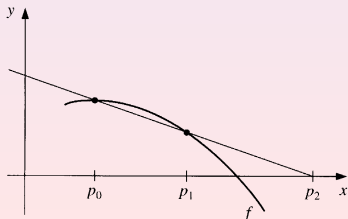


- Müller's method: Find complex roots of a polynomial  $P(x)$  (or any complex valued function  $f : \mathbb{C} \mapsto \mathbb{C}$ ):

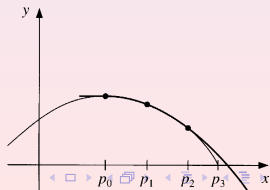
## Theorem

If  $z = a + ib$  is a complex zero of multiplicity  $m$  of  $P(x)$  with real coefficients, then  $\bar{z} = a - bi$  is also a zero of multiplicity  $m$  of  $P(x)$  and  $(x^2 - 2ax + a^2 + b^2)^m$  is a factor of  $P(x)$ .

**Secant method:** Given  $p_0$  and  $p_1$ , determine  $p_2$  as the intersection of the  $x$ -axis with the line through  $(p_0, f(p_0))$  and  $(p_1, f(p_1))$ .



**Müller's method:** Given  $p_0, p_1$  and  $p_2$ , determine  $p_3$  by the intersection of the  $x$ -axis with the parabola through  $(p_0, f(p_0))$ ,  $(p_1, f(p_1))$  and  $(p_2, f(p_2))$ .



Let

$$P(x) = a(x - p_2)^2 + b(x - p_2) + c$$

that passes through  $(p_0, f(p_0))$ ,  $(p_1, f(p_1))$  and  $(p_2, f(p_2))$ . Then

$$f(p_0) = a(p_0 - p_2)^2 + b(p_0 - p_2) + c,$$

$$f(p_1) = a(p_1 - p_2)^2 + b(p_1 - p_2) + c,$$

$$f(p_2) = a(p_2 - p_2)^2 + b(p_2 - p_2) + c = c.$$

It follows that

$$c = f(p_2),$$

$$b = \frac{(p_0 - p_2)^2 [f(p_1) - f(p_2)] - (p_1 - p_2)^2 [f(p_0) - f(p_2)]}{(p_0 - p_2)(p_1 - p_2)(p_0 - p_1)},$$

$$a = \frac{(p_1 - p_2) [f(p_0) - f(p_2)] - (p_0 - p_2) [f(p_1) - f(p_2)]}{(p_0 - p_2)(p_1 - p_2)(p_0 - p_1)}.$$



To determine  $p_3$ , a zero of  $P$ , we apply the quadratic formula to  $P(x) = 0$  and get

$$p_3 - p_2 = \frac{2c}{b \pm \sqrt{b^2 - 4ac}}. \quad (8)$$

If  $a, b, c$  are all real, we can choose

$$p_3 = p_2 + \frac{2c}{b + \operatorname{sgn}(b)\sqrt{b^2 - 4ac}}$$

such that the denominator will be largest in magnitude. The selected  $p_3$  is the one closer to  $p_2$  among those given in (8).

In case  $a, b, c$  are complex, the selection principle for  $p_3$  can be modified accordingly.





## Müller's method (Find a solution of $f(x) = 0$ )

Given  $p_0, p_1, p_2$ ; tolerance  $TOL$ ; maximum number of iterations  $M$

Set  $h_1 = p_1 - p_0$ ;  $h_2 = p_2 - p_1$ ;

$$\delta_1 = (f(p_1) - f(p_0))/h_1; \delta_2 = (f(p_2) - f(p_1))/h_2;$$

$$d = (\delta_2 - \delta_1)/(h_2 + h_1); i = 3.$$

While  $i \leq M$

$$\text{Set } b = \delta_2 + h_2 d; D = \sqrt{b^2 - 4f(p_2)d}.$$

If  $|b - D| < |b + D|$ , then set  $E = b + D$  else set  $E = b - D$ .

$$\text{Set } h = -2f(p_2)/E; p = p_2 + h.$$

If  $|h| < TOL$ , then STOP.

$$\text{Set } p_0 = p_1; p_1 = p_2; p_2 = p; h_1 = p_1 - p_0; h_2 = p_2 - p_1;$$

$$\delta_1 = (f(p_1) - f(p_0))/h_1; \delta_2 = (f(p_2) - f(p_1))/h_2;$$

$$d = (\delta_2 - \delta_1)/(h_2 + h_1); i = i + 1.$$

End while

